## Spectral Graph Sparsification

Presenter: Derrick Blakely<br>June 28, 2019<br>University of Virginia<br>https://qdata.github.io/deep2Read/

## Table of contents

1. Sparsification
2. The Graph Laplacian
3. Back to Sparsification
4. How to Obtain H
5. Sparsification using Effective Resistance

## Sparsification

## Intro

Approximate any graph with a sparse graph:


Approximation


## Why?

- Compression
- Faster to compute with
- Less memory


## Graph Approximation

Given: $G=(V, E, w)$ and $H=(V, F, z)$, we want:

$$
H \approx_{\epsilon} G
$$

Properties we want to preserve:

- Cut sizes
- Communities/clusters
- Behavior of random walks


## Cut Approximation

- $H \approx_{\epsilon} G$ if for every $S \subset V$, the sum of weights leaving $S$ is the same in $H$ and $G$



## Cut Approximation Theorem [2]

Given a graph $G$ with $m$ edges and error parameter $\epsilon$, we can find a graph $H$ such that:

- $H$ has $O\left(n \log n / \epsilon^{2}\right)$ edges
- The value of every cut in $H$ is $(1 \pm \epsilon)$ the corresponding cut in $G$
- $H$ can be constructed in $O\left(m \log ^{2} n\right)$ time if $G$ is unweighted and $O\left(m \log ^{3} n\right)$ if $G$ is weighted


## Stronger Approximation: Spectral Sparsification

$H \approx_{\kappa} G$ if for some error parameter $\kappa$ :

$$
H \preccurlyeq G \preccurlyeq \kappa H
$$

Where $H \preccurlyeq G$ if $\forall x: V \rightarrow \mathbb{R}$ :

$$
x^{\top} L_{H} x \leq x^{\top} L_{G} x
$$

## The Graph Laplacian

## The $\nabla$ Operator

- A pseudo-vector: $\nabla=\left[\frac{\partial}{\partial x_{1}}, \frac{\partial}{\partial x_{2}}, \ldots, \frac{\partial}{\partial x_{n}}\right]$


## The $\nabla$ Operator

- A pseudo-vector: $\nabla=\left[\frac{\partial}{\partial x_{1}}, \frac{\partial}{\partial x_{2}}, \ldots, \frac{\partial}{\partial x_{n}}\right]$
- Gradient: $\nabla f=\left[\frac{\partial f}{\partial x_{1}}, \frac{\partial f}{\partial x_{2}}, \ldots, \frac{\partial f}{\partial x_{n}}\right]$


## The $\nabla$ Operator

- A pseudo-vector: $\nabla=\left[\frac{\partial}{\partial x_{1}}, \frac{\partial}{\partial x_{2}}, \ldots, \frac{\partial}{\partial x_{n}}\right]$
- Gradient: $\nabla f=\left[\frac{\partial f}{\partial x_{1}}, \frac{\partial f}{\partial x_{2}}, \ldots, \frac{\partial f}{\partial x_{n}}\right]$
- Divergence: $\nabla \cdot f=\nabla \cdot\left[f_{1}, f_{2}, \ldots, f_{n}\right]=\frac{\partial f_{1}}{\partial x_{1}}+\frac{\partial f_{2}}{\partial x_{2}}+\ldots+\frac{\partial f_{n}}{\partial x_{n}}$


## The $\nabla$ Operator

- A pseudo-vector: $\nabla=\left[\frac{\partial}{\partial x_{1}}, \frac{\partial}{\partial x_{2}}, \ldots, \frac{\partial}{\partial x_{n}}\right]$
- Gradient: $\nabla f=\left[\frac{\partial f}{\partial x_{1}}, \frac{\partial f}{\partial x_{2}}, \ldots, \frac{\partial f}{\partial x_{n}}\right]$
- Divergence: $\nabla \cdot f=\nabla \cdot\left[f_{1}, f_{2}, \ldots, f_{n}\right]=\frac{\partial f_{1}}{\partial x_{1}}+\frac{\partial f_{2}}{\partial x_{2}}+\ldots+\frac{\partial f_{n}}{\partial x_{n}}$
- Laplacian: $\nabla \cdot \nabla f=\nabla^{2} f=\frac{\partial^{2} f}{\partial x_{1}^{2}}+\frac{\partial^{2} f}{\partial x_{2}^{2}}+\ldots+\frac{\partial^{2} f}{\partial x_{n}^{2}}$


## Physics Explanation

- Let $V: \mathbb{R}^{3} \rightarrow \mathbb{R}$ be the electric potential



## Physics Explanation

- Let $V: \mathbb{R}^{3} \rightarrow \mathbb{R}$ be the electric potential
- $E=-\nabla V: \mathbb{R}^{3} \rightarrow \mathbb{R}^{3}$ is the electric field

|  |
| :---: |
|  |  |
|  |  |
|  |  |
|  |  |

## Physics Explanation

－Let $V: \mathbb{R}^{3} \rightarrow \mathbb{R}$ be the electric potential
－$E=-\nabla V: \mathbb{R}^{3} \rightarrow \mathbb{R}^{3}$ is the electric field
－ $\operatorname{div}(E)=\nabla \cdot E: \mathbb{R}^{3} \rightarrow \mathbb{R}$ is divergence of $E$

| 11才フ－\11 |
| :---: |
| \11ハーマ 11 |
|  |
|  |
| 1 |

## Physics Explanation

－Let $V: \mathbb{R}^{3} \rightarrow \mathbb{R}$ be the electric potential
－$E=-\nabla V: \mathbb{R}^{3} \rightarrow \mathbb{R}^{3}$ is the electric field
－ $\operatorname{div}(E)=\nabla \cdot E: \mathbb{R}^{3} \rightarrow \mathbb{R}$ is divergence of $E$
－Laplacian $(V)=\operatorname{div}(E)=\nabla \cdot \nabla V=\nabla^{2} V$

| 11メーさ 11 |
| :---: |
| $11 /$ |
|  |
|  |
| t |

## Interpreting the Laplacian

- Laplacian $(V)=\operatorname{div}(E)=\nabla \cdot E=\nabla \cdot \nabla V=\nabla^{2} V$
- Extent to which a point behaves like a positive voltage source
- Net flux density through a volume at a point
- Second derivative of $V$ : smoothness of $V$ over space

| 11ヶ- \11 |
| :---: |
| 11入- - 1 |
|  |
|  |
|  |

## What is the Graph Laplacian?

How do we define $\nabla \cdot \nabla f=\nabla^{2} f$ for graphs? Three questions:

1. What does $f$ mean for graphs?
2. What does the gradient $\nabla f$ mean for graphs?
3. What does the Laplacian $\nabla \cdot \nabla f$ mean for graphs?

## 1. Functions Over Graphs

- $f: V \rightarrow \mathbb{R}$
- For example: the degree of each node
- In other words: degree of a node is like its potential



## 2. Gradient of Degree Function

- Need an analog to $\nabla=\left[\frac{\partial}{\partial x}, \frac{\partial}{\partial y}, \frac{\partial}{\partial z}\right]$
- Incidence matrix $K$ : each node gets a row and each edge gets a column
- If outgoing edge, $K_{n, e}=-1$
- If incoming edge, $K_{n, e}=1$
- Neither: $K_{n, e}=0$


## Incidence Matrix K



## Gradient of Degree Function



$$
\nabla f=K^{\top} f=K^{\top}\left[\begin{array}{l}
2  \tag{2}\\
2 \\
4 \\
1 \\
1
\end{array}\right]=\left[\begin{array}{c}
0 \\
2 \\
2 \\
-3 \\
-3
\end{array}\right]
$$

## 3. Laplacian Operator for Graphs

$$
\begin{align*}
\nabla \cdot \nabla=K K^{\top}=L & =\left[\begin{array}{ccccc}
2 & -1 & -1 & 0 & 0 \\
-1 & 2 & -1 & 0 & 0 \\
-1 & -1 & 4 & -1 & -1 \\
0 & 0 & -1 & 1 & 0 \\
0 & 0 & -1 & 0 & 1
\end{array}\right]  \tag{3}\\
L & =D-A \tag{4}
\end{align*}
$$

(where D is the degree matrix and A is the adjacency matrix)

## Graph Divergence



## Divergence and Smoothness



$$
\nabla^{2} f=K K^{\top} f=L f=\left[\begin{array}{c}
-2  \tag{6}\\
-2 \\
8 \\
-2 \\
-2
\end{array}\right]
$$

## Divergence and Smoothness



$$
\nabla^{2} f=K K^{\top} f=L f=\left[\begin{array}{c}
-1  \tag{7}\\
-1 \\
4 \\
-1 \\
-1
\end{array}\right]
$$

## Summary

| Physics | Graphs |
| :--- | ---: |
| Potential: $V$ | Vertex function: $f$ |
| $\nabla$ | $K:$ incidence matrix |
| $\nabla V:$ field | $K^{\top} f:$ graph gradient |
| $\nabla^{2}:$ Laplacian | $K K^{\top}$ Graph Laplacian |
| $\nabla^{2} V:$ Laplacian of $V$ | $K K^{\top} f: G r a p h$ Laplacian of $f$ |

## Back to Sparsification

## Stronger Approximation: Spectral Sparsification

$H \approx_{\epsilon} G$ if for some error parameter $\kappa$ :

$$
H \preccurlyeq G \preccurlyeq \kappa H
$$

Where $H \preccurlyeq G$ if $\forall x: V \rightarrow \mathbb{R}$ :

$$
x^{\top} L_{H} x \leq x^{\top} L_{G} x
$$

## Laplacian Quadratic Form

Sum of square differences across the edges:

$$
x^{\top} L_{G} x=\sum_{(u, v) \in E} w(u, v)(x(u)-x(v))^{2}
$$

## Spectral Approximation Theorem [1]

Given $G$ and error parameter $\epsilon$, we can find an approximation $H$ such that:

- $H \preccurlyeq G$
- $H$ has $O\left(n \log n / \epsilon^{2}\right)$ edges
- $H$ can be found in time $\tilde{O}\left(m / \epsilon^{2}\right)$


## Why is spectral approx. stronger than cut approx.?

- A spectral approximation is also a cut approximation
- The converse is not always true


## Why is it better?

- $H$ will inherit a bunch of properties of $G$
- The eigenvectors and eigenvalues will be similar
- Can use $H$ to obtain approximate solutions to linear systems of $G$
- If $x \in \mathbb{R}^{n}$, we can use convex solvers
- If $x \in\{0,1\}^{n}$, we can relax it to $x \in \mathbb{R}^{n}$ to obtain approximate solutions


## How to Obtain H

## Basic Idea: Random Sampling

- Choose edge $e$ with probability $p_{e}$
- Take $k$ independent samples
- Add e to $H$ with weight $1 / k p_{e}$


## Why not Uniform Sampling?

Don't want to disconnect the graph. E.g.,:


Instead: bias probabilities based on the "effective resistance" of the edges

# Sparsification using Effective Resistance 

## Resistor Networks



## Circuit Basics

Ohm's law: voltage drop across a resistor is

$$
V=I R
$$

Power dissipation across a resistor:

$$
P=V^{2} / R
$$

## Quadratic Form and Power Dissipation

- If we interpret $x$ as voltages and $E$ as conductances, then

$$
x^{\top} L_{G} x
$$

gives the power dissipation of the graph.

- If $H \approx_{\kappa} G$, then $H$ and $G$ have approximate "electrical equivalence."


## Effective Resistance

$R_{\text {eff }}(e)$ : power dissipation when a unit of current is sent across the ends


$$
R_{\text {eff }}(e)=\left\|L_{G}^{-1 / 2} b_{e}\right\|=b_{e}^{\top} L_{G}^{-1} b_{e}
$$

where $b_{e} \in 1,-1,0^{n}, b_{e}(u)=1, b_{e}(v)=-1,0$ everywhere else

## Sparsification with Effective Resistance [3]

- Choose edge $e$ with probability $p_{e} \propto R_{\text {eff }}(e)$
- Take $k$ independent samples
- Add $e$ to $H$ with weight $1 / k p_{e}$


## References

$\square$ J. Batson, D. A. Spielman, N. Srivastava, and S.-H. Teng.

Spectral sparsification of graphs: theory and algorithms.
Communications of the ACM, 56(8):87-94, 2013.
( A. Benczur and D. R. Karger.
Randomized approximation schemes for cuts and flows in capacitated graphs.
arXiv preprint cs/0207078, 2002.
戋 D. A. Spielman and N. Srivastava.
Graph sparsification by effective resistances.
SIAM Journal on Computing, 40(6):1913-1926, 2011.

