## **Spectral Graph Sparsification**

Presenter: Derrick Blakely

June 28, 2019

University of Virginia

https://qdata.github.io/deep2Read/

- 1. Sparsification
- 2. The Graph Laplacian
- 3. Back to Sparsification
- 4. How to Obtain H
- 5. Sparsification using Effective Resistance

## **Sparsification**

#### Approximate any graph with a sparse graph:



- Compression
- Faster to compute with
- Less memory

Given: 
$$G = (V, E, w)$$
 and  $H = (V, F, z)$ , we want:

 $H \approx_{\epsilon} G$ 

Properties we want to preserve:

- Cut sizes
- Communities/clusters
- Behavior of random walks

#### **Cut Approximation**

H ≈<sub>ε</sub> G if for every S ⊂ V, the sum of weights leaving S is the same in H and G



Given a graph G with m edges and error parameter  $\epsilon$ , we can find a graph H such that:

- *H* has  $O(n \log n/\epsilon^2)$  edges
- The value of every cut in H is  $(1\pm\epsilon)$  the corresponding cut in G
- H can be constructed in O(m log<sup>2</sup> n) time if G is unweighted and O(m log<sup>3</sup> n) if G is weighted

 $H \approx_{\kappa} G$  if for some error parameter  $\kappa$ :

 $H \preccurlyeq G \preccurlyeq \kappa H$ 

Where  $H \preccurlyeq G$  if  $\forall x : V \rightarrow \mathbb{R}$ :

 $x^{\top}L_{H}x \leq x^{\top}L_{G}x$ 

## The Graph Laplacian

• A pseudo-vector:  $\nabla = [\frac{\partial}{\partial x_1}, \frac{\partial}{\partial x_2}, ..., \frac{\partial}{\partial x_n}]$ 

• A pseudo-vector:  $\nabla = [\frac{\partial}{\partial x_1}, \frac{\partial}{\partial x_2}, ..., \frac{\partial}{\partial x_n}]$ 

• Gradient: 
$$\nabla f = \left[\frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2}, ..., \frac{\partial f}{\partial x_n}\right]$$

- A pseudo-vector:  $\nabla = [\frac{\partial}{\partial x_1}, \frac{\partial}{\partial x_2}, ..., \frac{\partial}{\partial x_n}]$
- Gradient:  $\nabla f = \left[\frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2}, ..., \frac{\partial f}{\partial x_n}\right]$
- Divergence:  $\nabla \cdot f = \nabla \cdot [f_1, f_2, ..., f_n] = \frac{\partial f_1}{\partial x_1} + \frac{\partial f_2}{\partial x_2} + ... + \frac{\partial f_n}{\partial x_n}$

• A pseudo-vector:  $\nabla = [\frac{\partial}{\partial x_1}, \frac{\partial}{\partial x_2}, ..., \frac{\partial}{\partial x_n}]$ 

• Gradient: 
$$\nabla f = \left[\frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2}, ..., \frac{\partial f}{\partial x_n}\right]$$

- Divergence:  $\nabla \cdot f = \nabla \cdot [f_1, f_2, ..., f_n] = \frac{\partial f_1}{\partial x_1} + \frac{\partial f_2}{\partial x_2} + ... + \frac{\partial f_n}{\partial x_n}$
- Laplacian:  $\nabla \cdot \nabla f = \nabla^2 f = \frac{\partial^2 f}{\partial x_1^2} + \frac{\partial^2 f}{\partial x_2^2} + \dots + \frac{\partial^2 f}{\partial x_n^2}$

• Let  $V: \mathbb{R}^3 \to \mathbb{R}$  be the electric potential



- Let  $V: \mathbb{R}^3 \to \mathbb{R}$  be the electric potential
- $E = -\nabla V : \mathbb{R}^3 \to \mathbb{R}^3$  is the electric field



- Let  $V: \mathbb{R}^3 
  ightarrow \mathbb{R}$  be the electric potential
- $E = -\nabla V : \mathbb{R}^3 \to \mathbb{R}^3$  is the electric field
- $div(E) = \nabla \cdot E : \mathbb{R}^3 \to \mathbb{R}$  is divergence of E



- Let  $V: \mathbb{R}^3 
  ightarrow \mathbb{R}$  be the electric potential
- $E = -\nabla V : \mathbb{R}^3 \to \mathbb{R}^3$  is the electric field
- $div(E) = \nabla \cdot E : \mathbb{R}^3 \to \mathbb{R}$  is divergence of E
- Laplacian(V) = div(E) =  $\nabla \cdot \nabla V = \nabla^2 V$



#### Interpreting the Laplacian

- Laplacian(V) = div(E) =  $\nabla \cdot E = \nabla \cdot \nabla V = \nabla^2 V$
- Extent to which a point behaves like a positive voltage source
- Net flux density through a volume at a point
- Second derivative of V: smoothness of V over space



How do we define  $\nabla \cdot \nabla f = \nabla^2 f$  for graphs? Three questions:

- 1. What does *f* mean for graphs?
- 2. What does the gradient  $\nabla f$  mean for graphs?
- 3. What does the Laplacian  $\nabla \cdot \nabla f$  mean for graphs?

#### 1. Functions Over Graphs

- $f: V \to \mathbb{R}$
- For example: the degree of each node
- In other words: degree of a node is like its potential



- Need an analog to  $\nabla = \left[\frac{\partial}{\partial x}, \frac{\partial}{\partial y}, \frac{\partial}{\partial z}\right]$
- Incidence matrix K: each node gets a row and each edge gets a column
- If outgoing edge,  $K_{n,e} = -1$
- If incoming edge,  $K_{n,e} = 1$
- Neither:  $K_{n,e} = 0$

#### Incidence Matrix K



(1)

#### **Gradient of Degree Function**



(2)

$$\nabla \cdot \nabla = KK^{\top} = L = \begin{bmatrix} 2 & -1 & -1 & 0 & 0 \\ -1 & 2 & -1 & 0 & 0 \\ -1 & -1 & 4 & -1 & -1 \\ 0 & 0 & -1 & 1 & 0 \\ 0 & 0 & -1 & 0 & 1 \end{bmatrix}$$
(3)  
$$L = D - A$$
(4)

(where D is the degree matrix and A is the adjacency matrix)

#### **Graph Divergence**



(5)

#### **Divergence and Smoothness**



(6)

#### **Divergence and Smoothness**



(7)

Physics	Graphs
Potential: V	Vertex function: $f$
$\nabla$	K: incidence matrix
$\nabla V$ : field	$K^{ op}f$ : graph gradient
$ abla^2$ : Laplacian	$KK^ op$ Graph Laplacian
$ abla^2 V$ : Laplacian of V	$KK^{\top}f$ : Graph Laplacian of $f$

## **Back to Sparsification**

 $H \approx_{\epsilon} G$  if for some error parameter  $\kappa$ :

 $H \preccurlyeq G \preccurlyeq \kappa H$ 

Where  $H \preccurlyeq G$  if  $\forall x : V \rightarrow \mathbb{R}$ :

 $x^{\top}L_{H}x \leq x^{\top}L_{G}x$ 

Sum of square differences across the edges:

$$x^{\top}L_G x = \sum_{(u,v)\in E} w(u,v)(x(u)-x(v))^2$$

Given G and error parameter  $\epsilon,$  we can find an approximation H such that:

- $H \preccurlyeq G$
- *H* has  $O(n \log n/\epsilon^2)$  edges
- *H* can be found in time  $ilde{O}(m/\epsilon^2)$

- A spectral approximation is also a cut approximation
- The converse is not always true

- H will inherit a bunch of properties of G
- The eigenvectors and eigenvalues will be similar
- Can use H to obtain approximate solutions to linear systems of G
- If  $x \in \mathbb{R}^n$ , we can use convex solvers
- If x ∈ {0,1}<sup>n</sup>, we can relax it to x ∈ ℝ<sup>n</sup> to obtain approximate solutions

## How to Obtain H

- Choose edge e with probability  $p_e$
- Take k independent samples
- Add e to H with weight  $1/kp_e$

#### Don't want to disconnect the graph. E.g.,:



Instead: bias probabilities based on the "effective resistance" of the edges

# Sparsification using Effective Resistance



Ohm's law: voltage drop across a resistor is

V = IR

Power dissipation across a resistor:

$$P = V^2/R$$

• If we interpret x as voltages and E as conductances, then

### $x^{\top}L_G x$

gives the power dissipation of the graph.

• If  $H \approx_{\kappa} G$ , then H and G have approximate "electrical equivalence."

 $R_{eff}(e)$ : power dissipation when a unit of current is sent across the ends ∽₩₩∽ R4 **R**5 **R**9 BC **R**2 of e  $R_{eff}(e) = ||L_{c}^{-1/2}b_{e}|| = b_{c}^{\top}L_{c}^{-1}b_{e}$ where  $b_e \in 1, -1, 0^n$ ,  $b_e(u) = 1$ ,  $b_e(v) = -1$ , 0 everywhere else

- Choose edge e with probability  $p_e \propto R_{eff}(e)$
- Take k independent samples
- Add e to H with weight  $1/kp_e$

- J. Batson, D. A. Spielman, N. Srivastava, and S.-H. Teng. Spectral sparsification of graphs: theory and algorithms. *Communications of the ACM*, 56(8):87–94, 2013.
- A. Benczur and D. R. Karger.
   Randomized approximation schemes for cuts and flows in capacitated graphs.

arXiv preprint cs/0207078, 2002.

D. A. Spielman and N. Srivastava. **Graph sparsification by effective resistances.** *SIAM Journal on Computing*, 40(6):1913–1926, 2011.